

Infinite Runs in Priced Timed Automata: Discounting

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Quasimodo 3

Motivation

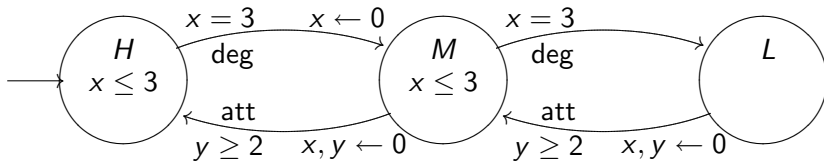
- **The overall goal:** Find optimal infinite paths in priced timed automata
- (Priced timed automaton; PTA: timed automaton with weights (“prices”) in locations and on the edges)
- Different versions:
 - Bouyer, Brinksma, Larsen, HSCC'04: *Staying alive as cheap as possible*: Given a PTA with two non-negative price functions, *cost* and *reward*, find infinite path with **lowest ratio accumulated cost/accumulated reward**
 - Bouyer, Fahrenberg, Larsen, Markey, Srba, FORMATS'08: *Infinite Runs in Weighted Timed Automata with Energy Constraints*: Given a PTA with positive or negative weights and possibly a pre-assigned threshold b , find infinite path for which accumulated cost ≥ 0 and possibly $\leq b$

Motivation

- This stuff: In the spirit of the HSCC'04 paper:
- One price function, all non-negative
- Interested in infinite paths with **lowest accumulated cost**
- But we apply **discounting**:
- Things which happen t time units in the future are taken into account only with a **discount** λ^t , for some fixed discounting factor λ
- (with **Kim Larsen**)
- INFINITY'08; QAPL'09

- 1 Motivation
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- 3 Proof
- 4 Summary
- 5 Fixed-point property

Priced timed automata

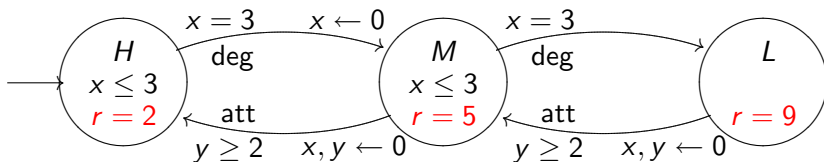


- Timed automaton:

Finite automaton

- + finite set of clocks $C = \{x, y, \dots\}$
- + location invariants
- + edge constraints
- + edge resets

Priced timed automata



- Timed automaton:

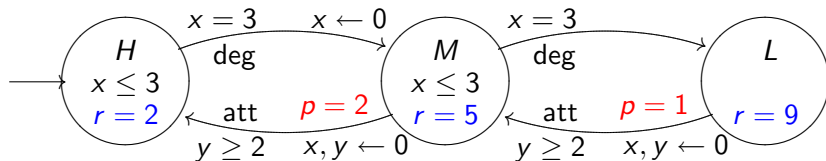
Finite automaton

- + finite set of clocks $C = \{x, y, \dots\}$
- + location invariants
- + edge constraints
- + edge resets

- Priced timed automaton:

- + price rates in locations (cost per time unit)

Priced timed automata



- Timed automaton:

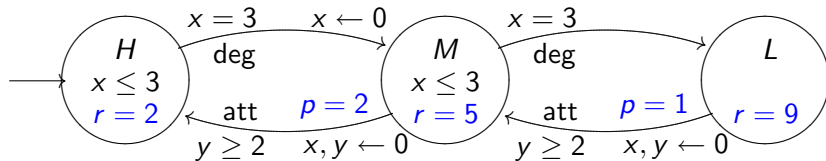
Finite automaton

- + finite set of clocks $C = \{x, y, \dots\}$
- + location invariants
- + edge constraints
- + edge resets

- Priced timed automaton:

- + price rates in locations (cost per time unit)
- + prices on transitions

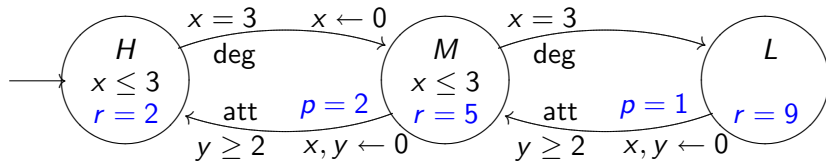
Priced timed automata



- An example (finite) run:

loc.	H	d	H	s	M	d	M	s	L	d	L	s	M	d	M	s	H
x	0		3		0		3		3		4		0		2		0
y	0		3		3		6		6		7		0		2		0
t	0		3		3		6		6		7		7		9		9

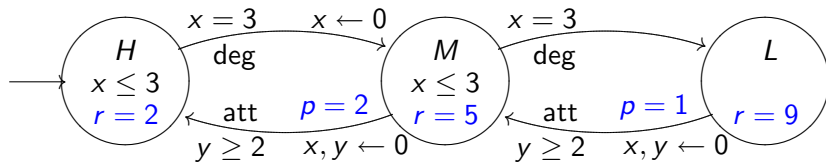
Priced timed automata



- An example (finite) run:

loc.	H	d	H	s	M	d	M	s	L	d	L	s	M	d
x	0		3		0		3		3		4		0	
y	0		3		3		6		6		7		0	
t	0		3		3		6		6		7		7	
	$2 \int_0^3 \lambda^t dt$		0	$5 \int_3^6 \lambda^t dt$		0	$9 \int_6^7 \lambda^t dt$		$1\lambda^7$	$5 \int_7^9 \lambda^t dt$				

Priced timed automata



- An example (finite) run:

loc.	H	d	H	s	M	d	M	s	L	d	L	s	M	d	M	s	H
x	0		3		0		3		3		4		0		2		0
y	0		3		3		6		6		7		0		2		0
t	0		3		3		6		6		7		7		9		9
$\lambda = .9$		5.14		0		9.38		0		4.54	.48		4.31		.77		

- Total discounted price of run: **24.62**
- Total discounted price of infinite loop: **40.2**

Discounted price

Formally: Let A be a **priced timed transition system** and $\lambda \in]0, 1]$.

- Discounted price of *finite* alternating path

$$\pi = s_0 \xrightarrow{t_0} s'_0 \rightarrow s_1 \rightarrow \cdots \xrightarrow{t_{n-1}} s'_{n-1} \rightarrow s_n :$$

$$P(\pi) = \sum_{i=0}^{n-1} \left(\int_{T_{i-1}}^{T_i} \lambda^t r(s_i^t) dt + \lambda^{T_i} p(s'_i \rightarrow s_{i+1}) \right)$$

with $T_i = \sum_{j=0}^i t_j$.

- Discounted price of *infinite* alternating path

$$\pi = s_0 \xrightarrow{t_0} s'_0 \rightarrow s_1 \rightarrow \cdots : \text{limit}$$

$$P(\pi) = \lim_{n \rightarrow \infty} P(s_0 \rightarrow s'_0 \rightarrow \cdots \rightarrow s'_{n-1} \rightarrow s_n)$$

provided that it exists. (!)

Problem and solution

Problem:

- Given: priced timed transition system A , state $s \in A$
- Find: an infinite path from s with **lowest discounted price**
- (or one that comes arbitrarily close)

Solution:

- For priced timed automata,
- which are *bounded*
- and *time-divergent*,
- and *rational* λ ,
- starting in the *initial state*,
- our problem is **computable**

Proof structure

Given priced timed automaton A :

- 1 Use **corner point abstraction** to construct *finite weighted graph* $\text{cp}(A)$ (refinement of *region graph*; Bouyer, Brinksma, Larsen, HSCC'04)
- 2 Find infinite path $\tilde{\pi}$ with lowest discounted price in $\text{cp}(A)$ using **linear programming** (*discounted payoff games*; Andersson, MSc. thesis, Uppsala '06)
- 3 Find cheapest path lying over $\tilde{\pi}$ in A

This needs a soundness and completeness result:

- Given an infinite path $\tilde{\pi}$ in $\text{cp}(A)$ for which $P(\tilde{\pi})$ converges, then for all $\varepsilon > 0$ there exists an infinite path $\pi \in \text{cp}^{-1}(\tilde{\pi})$ for which $|P(\pi) - P(\tilde{\pi})| < \varepsilon$.
- Given an infinite path π in A , there exists an infinite path $\tilde{\pi} \in \text{cp}(\pi)$ for which $P(\tilde{\pi}) \leq P(\pi)$.

Proof details

Given an infinite path π in A , there exists an infinite path $\tilde{\pi} \in \text{cp}(\pi)$ for which $P(\tilde{\pi}) \leq P(\pi)$.

Write $\pi = (q_0, \nu_0) \xrightarrow{p_0} (q_0, \nu_0 + t_0) \xrightarrow{p_1} (q_1, \nu_1) \xrightarrow{p_1} (q_1, \nu_1 + t_1) \xrightarrow{\dots} \dots$

Then

$$\begin{aligned} P(\pi) &= \sum_{i=0}^{\infty} \left(\int_{T_{i-1}}^{T_i} \lambda^t r(q_i) dt + \lambda^{T_i} p_i \right) \\ &= \sum_{i=0}^{\infty} \left(\frac{1}{\ln \lambda} r(q_i) (\lambda^{T_i} - \lambda^{T_{i-1}}) + p_i \lambda^{T_i} \right) \end{aligned}$$

– a function in variables T_0, T_1, \dots

\Rightarrow **Optimization problem:** Minimize $P(\pi) = f(T_0, T_1, \dots)$ under the constraint that T_0, T_1, \dots lie in a specific zone defined by π above.

Proof details, 2.

$$P(\pi) = f(T_0, T_1, \dots) = \sum_{i=0}^{\infty} \left(\frac{1}{\ln \lambda} r(q_i) (\lambda^{T_i} - \lambda^{T_{i-1}}) + p_i \lambda^{T_i} \right)$$

Task: Minimize $f(T_0, T_1, \dots)$ under the constraint that $(T_0, T_1, \dots) \in Z$ for a given (bounded, closed) zone Z .

-
- (can be shown that) f is (weakly) monotonic
 - easy to see: monotonic functions over *finite-dimensional* closed zones attain their minimum in a *corner point*
- ⇒ for *finite paths*, the corner point abstraction can “see” the path with lowest price (because it goes through corners) ⇒ done
- **Need:** Generalization of the above to **infinite-dimensional** zones
 - not easy, because infinite-dimensional zones are **not compact**
 - difficult part: show that infimum is attained **somewhere**

Proof details, 3.

Theorem: $Z \subseteq \mathbb{R}^\infty$ bounded and closed (in the supremum metric)

- f_1, f_2, \dots continuous functions $p_i : Z \rightarrow \mathbb{R}_{\geq 0}$ (*non-negative values!*)
- $f(x_1, x_2, \dots) = \sum_{i=1}^{\infty} f_i(x_i) : Z \rightarrow [0, \infty]$ converges for some $x \in Z$

\implies exists $z \in Z$ for which $f(z) = \inf_{y \in Z} f(y)$

Proof: Let $x : \mathbb{N} \rightarrow Z$ be a sequence for which $\lim f(x_i) = \inf_{y \in Z} f(y)$

Standard argument: Z is compact $\implies x$ contains converging subsequence $x' \implies$ let $z = \lim x' \implies$ done.

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But Z is not compact.

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$x = \quad x_1 \quad x_2 \quad x_3 \quad x_4 \quad x_5 \quad x_6 \quad x_7 \quad x_8 \quad x_9$

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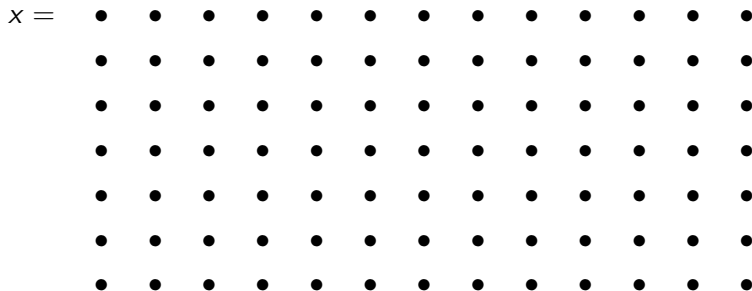
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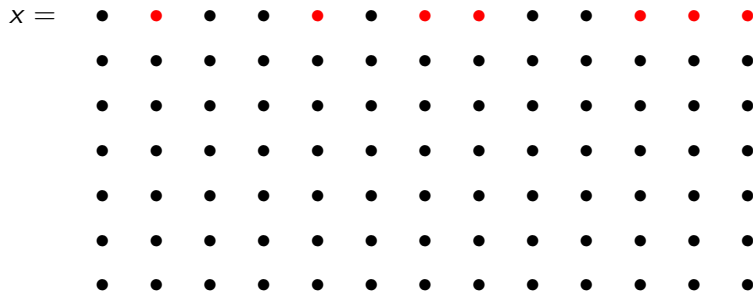
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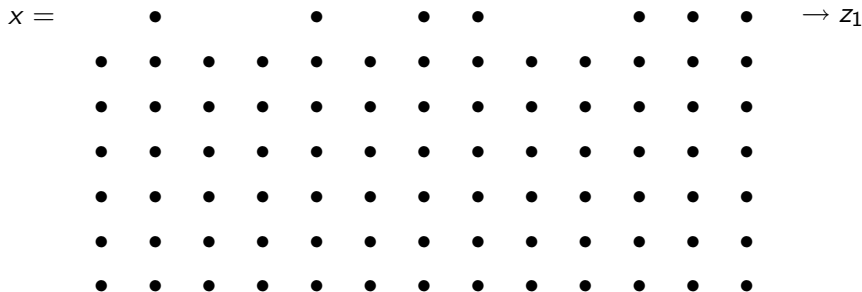
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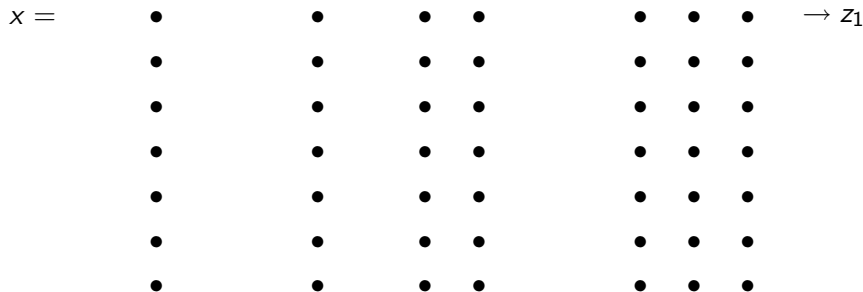
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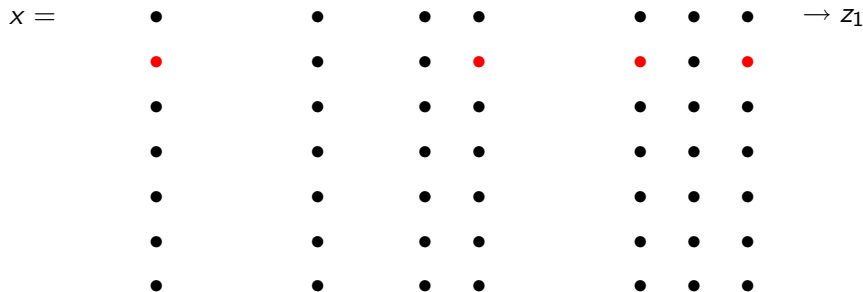
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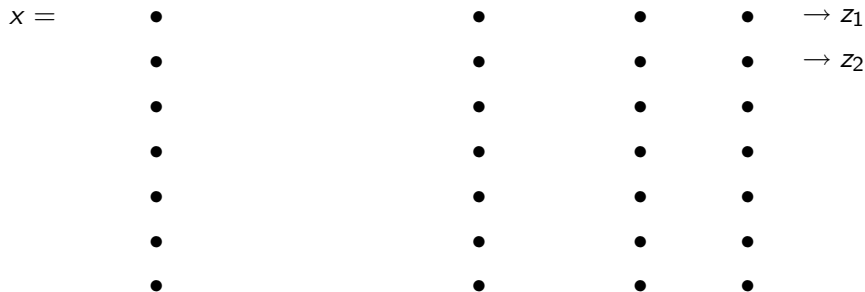
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- \implies exists $z \in Z$ for which $f(z) = \inf_{y \in Z} f(y)$

Corollary: $Z \subseteq \mathbb{R}^\infty$ bounded and closed **zone**

- f_1, f_2, \dots **monotonous** continuous functions $\text{pr}_i Z \rightarrow \mathbb{R}_{\geq 0}$
 - $f(x_1, x_2, \dots) = \sum_{i=1}^{\infty} f_i(x_i) : Z \rightarrow [0, \infty]$ converges for some $x \in Z$
- \implies exists **corner point** z of Z for which $f(z) = \inf_{y \in Z} f(y)$

So corner point abstraction *can* see the path with lowest price \implies done

Summary

- The discount-optimal infinite path problem is **computable for priced timed automata**
 - (with arbitrarily many clocks)
 - (under certain restrictions; only “real” restriction: time divergence)
 - (time-divergence assumption similar to reward-divergence assumption in **HSCC'04** paper; can it be avoided?)
 - Uses generalization of a well-known fact about monotonous functions defined on zones, to **infinite zones**
 - (Other applications?)
-

- Computable **yes**, feasible **NO**
- For zone-based iterative computations: need some **fixed-point property** / **additivity**

Additivity

- For \rightarrow a switch and π a path out of the target of the switch:

$$P(\rightarrow \circ \pi) = p(\rightarrow) + P(\pi)$$

For \xrightarrow{t} a delay and π a path out of the end state of the delay:

$$P(\xrightarrow{t} \circ \pi) = p(\xrightarrow{t}) + \lambda^t P(\pi)$$

- Similar property not satisfied for other formalisms (e.g. cost / reward, or step-wise discounting)
- Question: What do we know about the cost function if we require

$$P(\xrightarrow{t} \circ \pi) = p(\xrightarrow{t}) + g(t)P(\pi)$$

- (Natural property also because of **additivity of delays!**)
- Answer: **Everything!**

Additivity

Theorem: If P is an accumulated price function in a PTA which satisfies

$$P(s \xrightarrow{t_1} s' \xrightarrow{t_2} s'') = P(s \xrightarrow{t_1} s') + g(t_1) P(s' \xrightarrow{t_2} s'')$$

for all states and delays, then

$$g(t) = \lambda^t \quad \text{and} \quad P(s \xrightarrow{t} s') = \alpha(s) \int_0^t \lambda^t dt$$

for some $\lambda \in \mathbb{R}_{\geq 0}$ and $\alpha : S \rightarrow \mathbb{R}$.

Reason: The functional equation

$$f(t_1 + t_2) = f(t_1) + g(t_1) f(t_2)$$

has essentially only one solution.

Proof: Just do like Cauchy in 1821 !

Conclusion

- If you want to discount in real-time formalisms, use our kind of discounting
- (unless you have your own good reasons not to; but you'll lose additivity)
- Then you'll get computability of optimal infinite paths
- And if you give us a little more time, maybe you'll even have an efficient zone-based algorithm

