*-Continuous Kleene ω-Algebras
Theory and Applications (no Tools)

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WATA 2016

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History

- WATA 2012: *Büchi conditions for generalized energy automata*
- ATVA 2013: *Kleene algebras and semimodules for energy problems*
- DLT 2015: **-continuous Kleene ω-algebras*
- FICS 2015: **-continuous Kleene ω-algebras for energy problems*

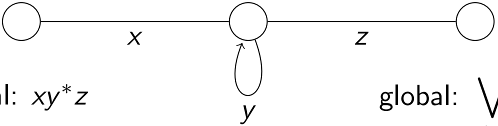
History

- FORMATS 2008: *Infinite runs in weighted timed automata with energy constraints*
- HSCC 2010: *Timed automata with observers under energy constraints*
- WATA 2012: *Büchi conditions for generalized energy automata*
- ATVA 2013: *Kleene algebras and semimodules for energy problems*
- DLT 2015: **-continuous Kleene ω-algebras*
- FICS 2015: **-continuous Kleene ω-algebras for energy problems*

*-Continuous Kleene Algebras

- idempotent semiring $S = (S, \vee, \cdot, \bot, 1)$
- in which all infinite suprema $x^* := \bigvee\{x^n \mid n \geq 0\}$ exist,
- and such that for all $x, y, z \in S$, $xy^*z = \bigvee_{n\geq 0} xy^n z$

Consequence: loop abstraction:



local: xy^*z

global: $\bigvee_{n\geq 0} xy^n z$

Continuous Kleene Algebras

- Kleene algebra $S = (S, \vee, \cdot, ^*, \bot, 1)$
- in which all suprema $\bigvee X$, $X \subseteq S$ exist,
- and such that for all $X \subseteq S$, $y, z \in S$, $y(\bigvee X)z = \bigvee yXz$

- All continuous Kleene algebras are *-continuous, but not vice-versa
  - Example: regular languages over some Σ

Matrix Semirings

S semiring, $n \geq 1$

- $S^{n\times n}$: semiring of $n \times n$-matrices over S
- (with matrix addition and multiplication)
- If S is a *-continuous Kleene algebra, then so is $S^{n\times n}$
- with $M^*_{i,j} = \bigvee_{m\geq 0} \bigvee_{1\leq k_1,\dots,k_m\leq n} M_{i,k_1} M_{k_1,k_2} \cdots M_{k_m,j}$
- and for $M = \begin{bmatrix} a & b \\ c & d \end{bmatrix}$,

$$M^* = \begin{bmatrix} (a \vee bd^*c)^* & (a \vee bd^*c)^* bd^* \\ (d \vee ca^*b)^* ca^* & (d \vee ca^*b)^* \end{bmatrix}$$

(recursively)

Finite Runs in Weighted Automata

S *-continuous Kleene algebra, $n \geq 1$

- a weighted automaton over S (with n states): $A = (\alpha, M, \kappa)$
- $\alpha \in \{\bot, 1\}^n$ initial vector, $\kappa \in \{\bot, 1\}^n$ accepting vector, $M \in S^{n \times n}$ transition matrix
- finite behavior of A: $|A| = \alpha M^* \kappa$
- Theorem:

$$|A| = \bigvee \left\{ w_0 \cdots w_n \;\middle|\; s_i \xrightarrow{w_0} \cdots \xrightarrow{w_n} s_j \text{ accepting path in } S \right\}$$

$(s_i \xrightarrow{w_0} \cdots \xrightarrow{w_n} s_j$ accepting if $\alpha_i = \kappa_j = 1)$

Idempotent Semiring-Semimodule Pairs

- idempotent semiring $S = (S, \vee, \cdot, \bot, 1)$
- commutative idempotent monoid $V = (V, \vee, \bot)$
- left S-action $S \times V \to V$, $(s, v) \mapsto sv$
- such that for all $s, s' \in S$, $v \in V$:

$$(s \vee s')v = sv \vee s'v \qquad s(v \vee v') = sv \vee sv'$$
$$(ss')v = s(s'v) \qquad \bot s = \bot$$
$$s\bot = \bot \qquad 1v = v$$

Continuous Kleene ω-Algebras

- idempotent semiring-semimodule pair (S, V)
- where S is a continuous Kleene algebra,
- V is a complete lattice,
- and the S-action on V preserves all suprema in either argument,
- with an infinite product $\prod : S^\omega \to V$ such that:
  - For all $x_0, x_1, \ldots \in S$, $\prod x_n = x_0 \prod x_{n+1}$.
  - Let $x_0, x_1, \ldots \in S$ and $0 = n_0 \le n_1 \le \cdots$ a sequence which increases without a bound. Let $y_k = x_{n_k} \cdots x_{n_{k+1}-1}$ for all $k \ge 0$. Then $\prod x_n = \prod y_k$.
  - For all $X_0, X_1, \ldots \subseteq S$, $\prod(\bigvee X_n) = \bigvee\{\prod x_n \mid x_n \in X_n, n \ge 0\}$.

Matrix Semiring-Semimodule Pairs

(S, V) semiring-semimodule pair, $n \geq 1$

- $(S^{n \times n}, V^n)$ is again a semiring-semimodule pair
- (the action is matrix-vector product)
- If (S, V) is a continuous Kleene ω-algebra, then so is $(S^{n \times n}, V^n)$
- with $M_i^\omega = \bigvee_{1 \leq k_1, k_2, \ldots \leq n} M_{i,k_1} M_{k_1,k_2} \cdots$
- and for $M = \begin{bmatrix} a & b \\ c & d \end{bmatrix}$,

$$M^\omega = \begin{bmatrix} (a \vee bd^*c)^\omega \vee (a \vee bd^*c)^* bd^\omega \\ (d \vee ca^*b)^\omega \vee (d \vee ca^*b)^* ca^\omega \end{bmatrix}$$

(recursively)

Infinite Runs in Weighted Automata

(S, V) continuous Kleene ω-algebra
(α, M, κ) weighted automaton over S

- Reorder $S = \{1, \dots, n\}$ so that $\kappa = (1, \dots, 1, \bot, \dots, \bot)$
  - i.e. the first $k \le n$ states are accepting
- Büchi behavior of A: write $M = \begin{bmatrix} a & b \\ c & d \end{bmatrix}$, with $a \in S^{k \times k}$, then

$$\|A\| = \alpha \begin{bmatrix} (a + bd^* c)^\omega \\ d^* c(a + bd^* c)^\omega \end{bmatrix}$$

- Theorem:

$$\|A\| = \bigvee \Big\{ \prod w_n \;\Big|\; s_i \xrightarrow{w_0} \xrightarrow{w_1} \cdots \text{ Büchi path in } S \Big\}$$

($s_i \xrightarrow{w_0} \xrightarrow{w_1} \cdots$ Büchi path if $\alpha_i = 1$ and some s_j with $j \le k$ is visited infinitely often)

Problem

continuous Kleene algebras	continuous Kleene ω-algebras
*-continuous Kleene algebras	???

Problem

continuous Kleene algebras	continuous Kleene ω-algebras
*-continuous Kleene algebras	*-continuous Kleene ω-algebras

[Esik, F., Legay 2015 (DLT)]

Generalized *-Continuous Kleene Algebras [EFL'15]

- semiring-semimodule pair (S, V)
- where S is a *-continuous Kleene algebra
- such that for all $x, y \in S$, $v \in V$, $xy^*v = \bigvee_{n \geq 0} xy^n v$

*-Continuous Kleene ω-Algebras [EFL'15]

- generalized *-continuous Kleene algebra (S, V)
- with an infinite product $\prod : S^\omega \to V$ such that:
  - For all $x_0, x_1, \ldots \in S$, $\prod x_n = x_0 \prod x_{n+1}$.
  - Let $x_0, x_1, \ldots \in S$ and $0 = n_0 \le n_1 \le \cdots$ a sequence which increases without a bound. Let $y_k = x_{n_k} \cdots x_{n_{k+1}-1}$ for all $k \ge 0$. Then $\prod x_n = \prod y_k$.
  - For all $x_0, x_1, \ldots, y, z \in S$, $\prod (x_n(y \vee z)) = \bigvee_{x'_0, x'_1, \ldots \in \{y, z\}} \prod x_n x'_n$.
  - For all $x, y_0, y_1, \ldots \in S$, $\prod x^* y_n = \bigvee_{k_0, k_1, \ldots \ge 0} \prod x^{k_n} y_n$.

Matrix Semiring-Semimodule Pairs, Revisited [EFL'15]

(S, V) *-continuous Kleene ω-algebra, $n \geq 1$

- $(S^{n\times n}, V^n)$ is a generalized *-continuous Kleene algebra
- with an operation $^\omega : S^{n\times n} \to V^n$ given by
  $M_i^\omega = \bigvee_{1\leq k_1, k_2, \ldots \leq n} M_{i,k_1} M_{k_1,k_2} \cdots$
  - (not a general infinite product)
- and for $M = \begin{bmatrix} a & b \\ c & d \end{bmatrix}$,

$$M^\omega = \begin{bmatrix} (a \vee bd^*c)^\omega \vee (a \vee bd^*c)^* bd^\omega \\ (d \vee ca^*b)^\omega \vee (d \vee ca^*b)^* ca^\omega \end{bmatrix}$$

(recursively)

Free Continuous Kleene ω-Algebras

Let A be a set.

Theorem (old)

The language semiring $P(A^*)$ is the free continuous Kleene algebra on A.

Theorem

$(P(A^*), P(A^\infty))$ *is the free continuous Kleene ω-algebra on A.*

Theorem

$(P(A^*), P(A^\omega))$ *is the free continuous Kleene ω-algebra on A satisfying $1^\omega = \bot$.*

Free Finitary *-Continuous Kleene ω-Algebras

Let A be a set.

Theorem (old)

*The regular language semiring $R(A^*)$ is the free *-continuous Kleene algebra on A.*

Theorem

$(R(A^*), R'(A^\infty))$ *is the free finitary *-continuous Kleene ω-algebra on A.*

Theorem

$(R(A^*), R'(A^\omega))$ *is the free finitary *-continuous Kleene ω-algebra satisfying $1^\omega = \perp$ on A.*

- finitary: ω-product only defined for finitary sequences
- $R'(A^\omega)$: finite unions of finitary infinite products
- know nothing about the non-finitary case

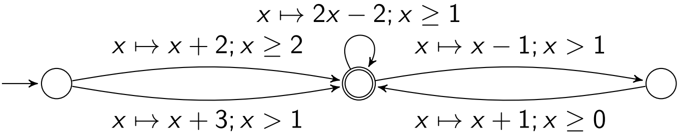
Energy Automata

Energy function:

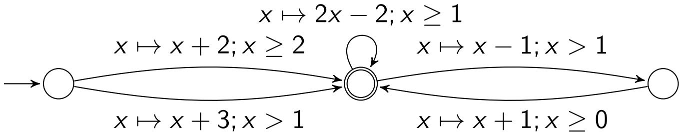
- partial function $f : \mathbb{R}_{\ge 0} \hookrightarrow \mathbb{R}_{\ge 0}$
- which is defined on some closed interval $[l_f, \infty[$ or on some open interval $]l_f, \infty[$,
- and such that for all $x \le y$ for which f is defined,

$$f(y) - f(x) \ge y - x$$

Energy automaton: finite automaton labeled with energy functions



Energy Automata, Semantically



- Start with initial energy x_0 and update at transitions according to label function
- If label function undefined on input, transition is disabled

Reachability: Given x_0, does there exist an accepting (finite) run with initial energy x_0?

Büchi: Given x_0, does there exist a Büchi (infinite) run with initial energy x_0?

Energy Automata, Algebraically

- Let $L = [0, \infty]_\perp$: extended nonnegative real numbers plus bottom
  - (a complete lattice)
- Extended energy function: function $f : L \to L$
- with $f(\perp) = \perp$, and $f(\infty) = \infty$ unless $f(x) = \perp$ for all $x \in L$,
- and $f(y) - f(x) \geq y - x$ for all $x \leq y$.
- Set $\mathcal{E}$ of such functions is an idempotent semiring with operations $\vee$ (pointwise max) and $\circ$ (composition)
- in fact, a *-continuous Kleene algebra
  - $f^*(x) = x$ if $f(x) \leq x$; $f^*(x) = \infty$ if $f(x) > x$
  - not a continuous Kleene algebra

Theorem (Reachability)

There exists an accepting (finite) run from initial energy x_0 iff $|A|(x_0) \neq \perp$.

Energy Automata, Algebraically, 2.

- Let $\mathbf{2} = \{\mathbf{ff}, \mathbf{tt}\}$: the Boolean lattice
- Let $\mathcal{V}$ be the set of monotone and $\top$-continuous functions $L \to \mathbf{2}$
  - $f : L \to \mathbf{2}$ $\top$-continuous if $f(x) \equiv \mathbf{ff}$ or for all $X \subseteq L$ with $\bigvee X = \infty$, also $\bigvee f(X) = \mathbf{tt}$.
- $(\mathcal{E}, \mathcal{V})$ is an idempotent semiring-semimodule pair
- Define $\prod : \mathcal{E}^\omega \to \mathcal{V}$ by

$$(\prod f_n)(x) = \mathbf{tt} \text{ iff } \forall n \geq 0 : f_n(f_{n-1}(\cdots(x)\cdots)) \neq \bot$$

- Lemma: $\prod f_n$ is indeed $\top$-continuous for all $f_0, f_1, \ldots \in \mathcal{E}$
- Theorem: $(\mathcal{V}, \mathcal{E})$ is a *-continuous Kleene ω-algebra
  - not a continuous Kleene ω-algebra

Theorem

Büchi There exists a Büchi run from initial energy x_0 iff $\|A\|(x_0) \neq \mathbf{ff}$.

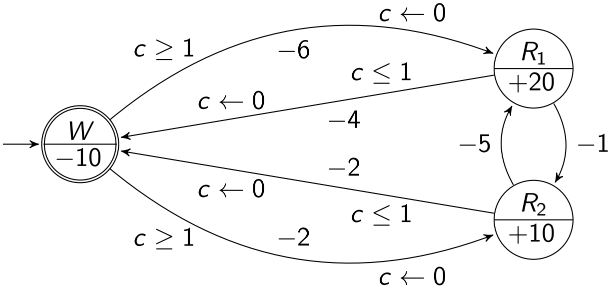
Conclusion

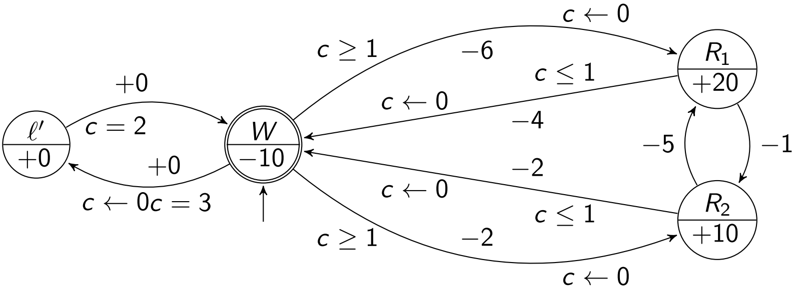
- *-continuous Kleene ω-algebras: a useful generalization of continuous Kleene ω-algebras
  - (like *-continuous Kleene algebras are a useful generalization of continuous Kleene algebras)
- can be used to solve general one-dimensional energy problems

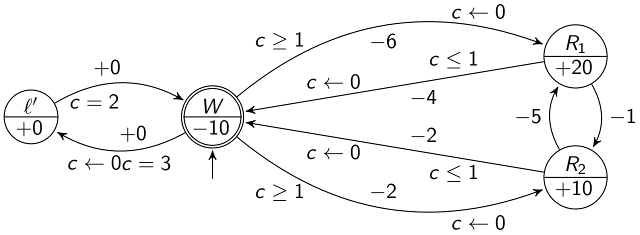
Other work:

- real-time energy problems (FSTTCS 2015)
- higher-dimensional energy problems?
- hybrid systems?
- non-idempotent case?

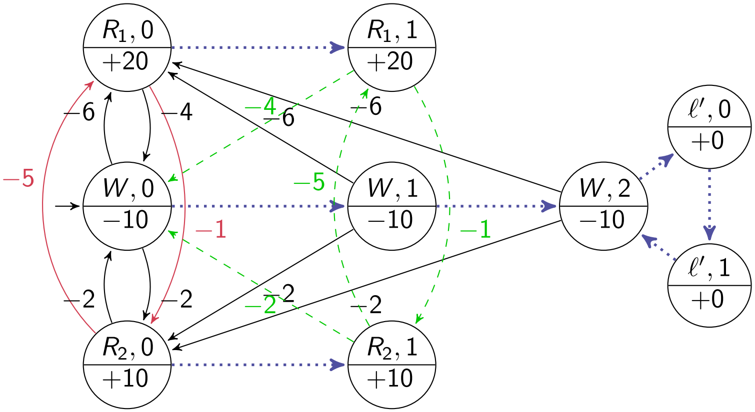
5 From Timed Energy Problems to Energy Automata


W
−10
R1
+20
R2
+10
c ≥ 1
−6
c ← 0
c ≤ 1
c ← 0
−4
−5
−1
−2
c ← 0
c ≤ 1
c ≥ 1
−2
c ← 0

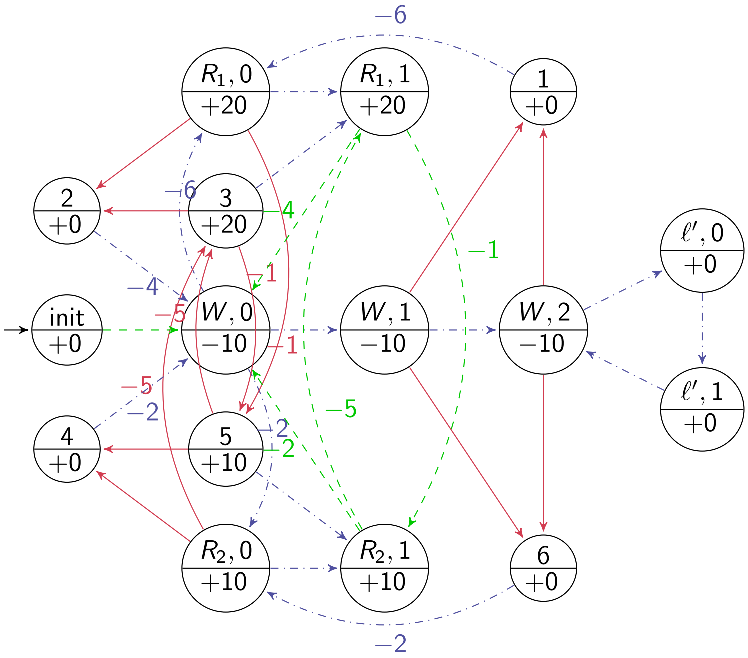

ℓ′
+0
W
−10
R1
+20
R2
+10
+0
c = 2
+0
c ← 0c = 3
c ≥ 1
−6
c ← 0
c ≤ 1
c ← 0
−4
−5
−1
−2
c ← 0
c ≤ 1
c ≥ 1
−2
c ← 0



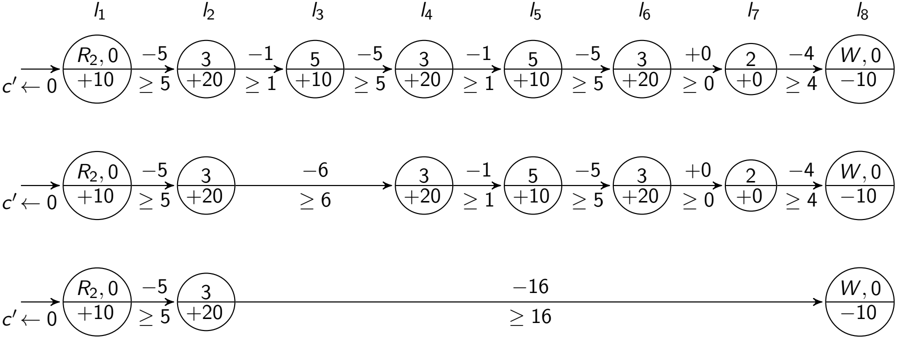
ℓ′
+0
W
−10
R1
+20
R2
+10
+0
c = 2
+0
c ← 0c = 3
c ≥ 1
−6
c ← 0
c ≤ 1
c ← 0
−4
−5
−1
−2
c ← 0
c ≤ 1
c ≥ 1
−2
c ← 0



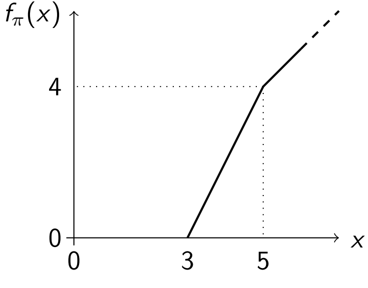
R1, 0
+20
R1, 1
+20
W, 0
−10
W, 1
−10
W, 2
−10
R2, 0
+10
R2, 1
+10
ℓ′, 0
+0
ℓ′, 1
+0
−6
−4
−4
−6
−6
−5
−5
−1
−1
−2
−2
−2
−2
−2



−6
R1, 0
+20
R1, 1
+20
1
+0
2
+0
3
+20
−6
−4
−1
−1
ℓ′, 0
+0
−4
init
+0
−5
W, 0
−10
−1
W, 1
−10
W, 2
−10
ℓ′, 1
+0
−5
−2
−5
4
+0
5
+10
−2
−2
R2, 0
+10
R2, 1
+10
6
+0
−2



l_1 l_2 l_3 l_4 l_5 l_6 l_7 l_8
$c' \leftarrow 0$
$R_2, 0$ / $+10$
-5 / ≥ 5
3 / $+20$
-1 / ≥ 1
5 / $+10$
-5 / ≥ 5
3 / $+20$
-1 / ≥ 1
5 / $+10$
-5 / ≥ 5
3 / $+20$
$+0$ / ≥ 0
2 / $+0$
-4 / ≥ 4
$W, 0$ / -10
-6 / ≥ 6
-16 / ≥ 16



$f_\pi(x)$
4
0
x
0
3
5