Quantitative Verification

The Good, The Bad and The Ugly

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Nice People

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Model Checking

model		specification
Mod	$\models$	Spec

Quantitative Model Checking

quantitative model		quantitative specification
Mod	$\models$	Spec

Quantitative Model Checking

quantitative model | quantitative specification

$$\text{Mod} \models \text{Spec}$$

↑

not sufficient

replace by

$$\models_\varepsilon$$

Claus T: Quantitative Quantitative Quantitative Analysis

Quantitative *Models*

Quantitative *Logics*

$$\Pr_{\leq .1}(\Diamond \mathit{error})$$

Quantitative *Verification*

$$[\![\phi]\!](s) = 3.14$$
$$d(s,t) = 42$$

Boolean world	**"Quantification"**
Trace equivalence $\equiv$	Linear distances d_L
Bisimilarity $\sim$	Branching distances d_B
$s \sim t$ implies $s \equiv t$	$d_L(s,t) \leq d_B(s,t)$
$s \models \phi$ or $s \not\models \phi$	$[\![\phi]\!](s)$ is a quantity
$s \sim t$ iff $\forall \phi : s \models \phi \Leftrightarrow t \models \phi$	$d_B(s,t) = \sup_\phi d([\![\phi]\!](s), [\![\phi]\!](t))$

Compositional Verification

model specification

$$\text{Mod} \models \text{Spec}$$

- $\text{Mod} \models \text{Spec}_1$ & $\text{Spec}_1 \leq \text{Spec}_2 \implies \text{Mod} \models \text{Spec}_2$
- $\text{Mod} \models \text{Spec}_1$ & $\text{Mod} \models \text{Spec}_2 \implies \text{Mod} \models \text{Spec}_1 \wedge \text{Spec}_2$
- $\text{Mod}_1 \models \text{Spec}_1$ & $\text{Mod}_2 \models \text{Spec}_2 \implies \text{Mod}_1 \| \text{Mod}_2 \models \text{Spec}_1 \| \text{Spec}_2$
- $\text{Mod}_1 \models \text{Spec}_1$ & $\text{Mod}_2 \models \text{Spec}/\text{Spec}_1 \implies \text{Mod}_1 \| \text{Mod}_2 \models \text{Spec}$

- bottom-up and top-down

Quantitative Compositional Verification?

quantitative model — quantitative specification

$$\text{Mod} \models_\varepsilon \text{Spec}$$

- $\text{Mod} \models_\varepsilon \text{Spec}_1$ & $\text{Spec}_1 \le_\varepsilon \text{Spec}_2 \implies \text{Mod} \models_\varepsilon \text{Spec}_2$
- $\text{Mod} \models_\varepsilon \text{Spec}_1$ & $\text{Mod} \models_\varepsilon \text{Spec}_2 \implies \text{Mod} \models_\varepsilon \text{Spec}_1 \wedge \text{Spec}_2$
- $\text{Mod}_1 \models_\varepsilon \text{Spec}_1$ & $\text{Mod}_2 \models_\varepsilon \text{Spec}_2 \implies \text{Mod}_1 \| \text{Mod}_2 \models_\varepsilon \text{Spec}_1 \| \text{Spec}_2$
- $\text{Mod}_1 \models_\varepsilon \text{Spec}_1$ & $\text{Mod}_2 \models_\varepsilon \text{Spec}/\text{Spec}_1 \implies \text{Mod}_1 \| \text{Mod}_2 \models_\varepsilon \text{Spec}$

- surely not the same ε everywhere!?

User Stories

"In your quantitative verification, what type of distances do you use?"

- point-wise $D(\sigma,\tau) = \sup_i |\sigma_i - \tau_i|$
- accumulating $D(\sigma,\tau) = \sum_i |\sigma_i - \tau_i|$
- limit-average $D(\sigma,\tau) = \limsup_N \frac{1}{N} \sum_{i=0}^N |\sigma_i - \tau_i|$
- discounted $D(\sigma,\tau) = \sum_i \lambda^i |\sigma_i - \tau_i|$
- maximum-lead $D(\sigma,\tau) = \sup_N \left| \sum_{i=0}^N (\sigma_i - \tau_i) \right|$
- Cantor $D(\sigma,\tau) = 1/(1 + \inf\{j \mid \sigma_j \neq \tau_j\})$
- discrete $D(\sigma,\tau) = 0$ if $\sigma = \tau$; ∞ otherwise

Asarin-Basset-Degorre 2018

$$D(\sigma, \tau) = \max \begin{cases} \sup_i \inf_j \{|t_i - s_j| \mid a_i = b_j\} \\ \sup_j \inf_i \{|t_i - s_j| \mid a_i = b_j\} \end{cases}$$

Challenge (*ca.* 2009)

- In quantitative verification, lots of different distances
- Develop theory to cover all/most of them
  - idea: use bisimulation games

⇒ The Quantitative Linear-Time–Branching-Time Spectrum

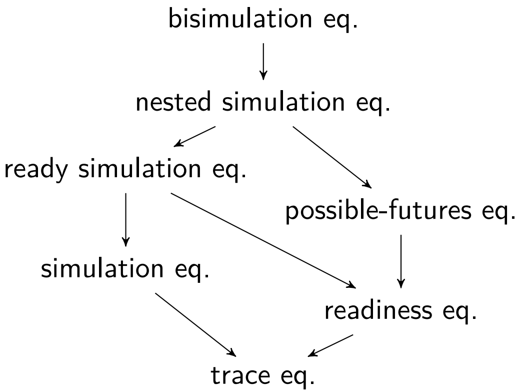
  - QAPL 2011, FSTTCS 2011, TCS 2014

Challenge (*ca.* 2012):

- How to make this compositional?
- Still not satisfied!

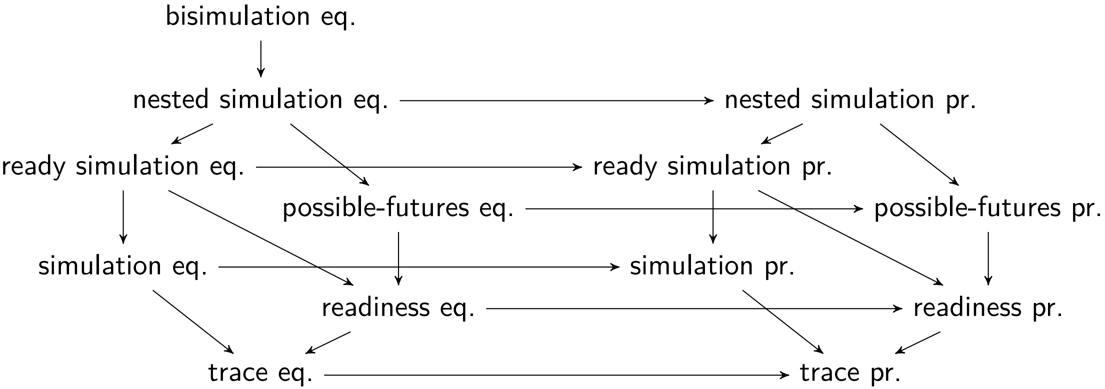
The Linear-Time–Branching-Time Spectrum

van Glabbeek 1990 (excerpt):



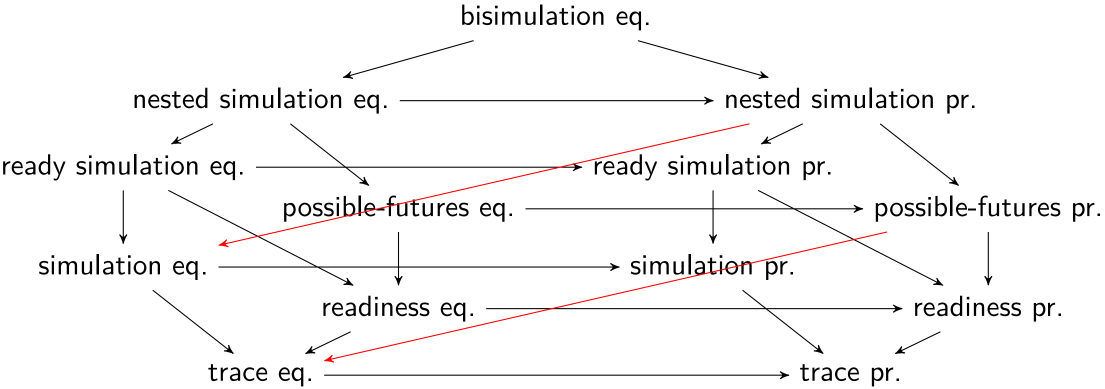
The Linear-Time–Branching-Time Spectrum

van Glabbeek 1990 (excerpt):

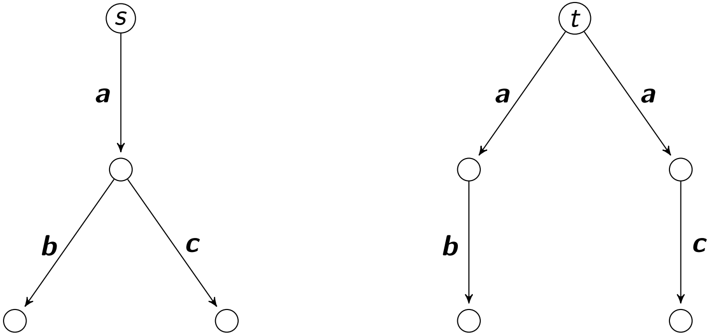


The Linear-Time–Branching-Time Spectrum

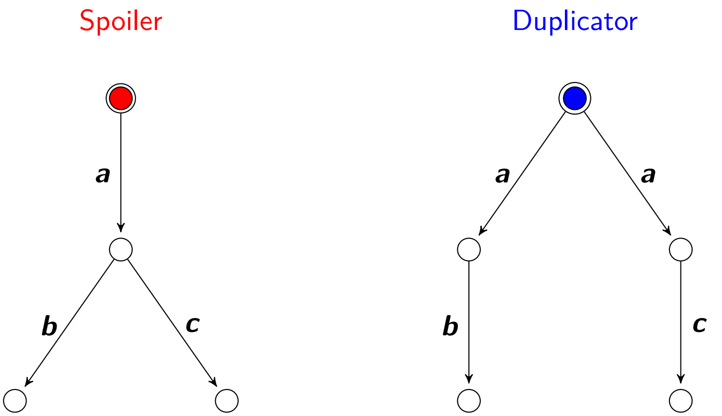
van Glabbeek 1990 (excerpt):



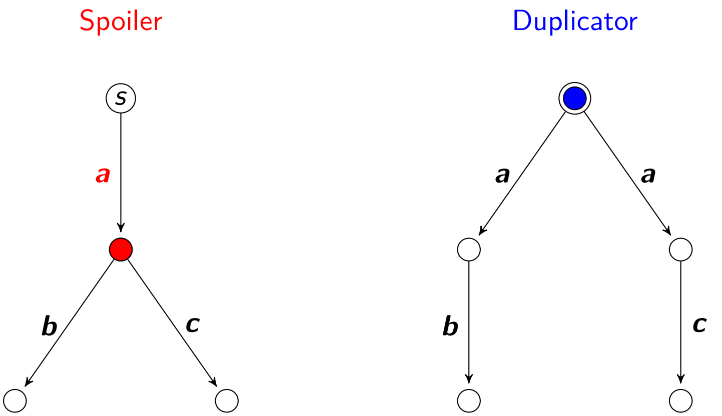
The Simulation Game



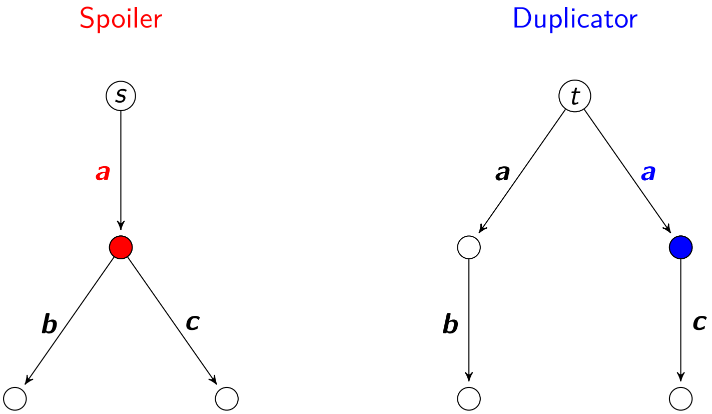
The Simulation Game



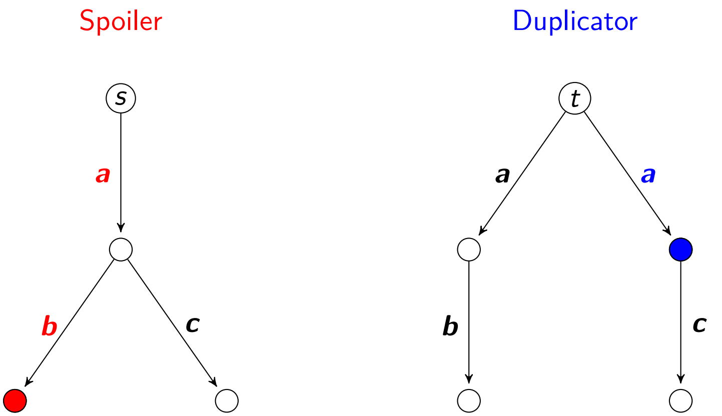
The Simulation Game



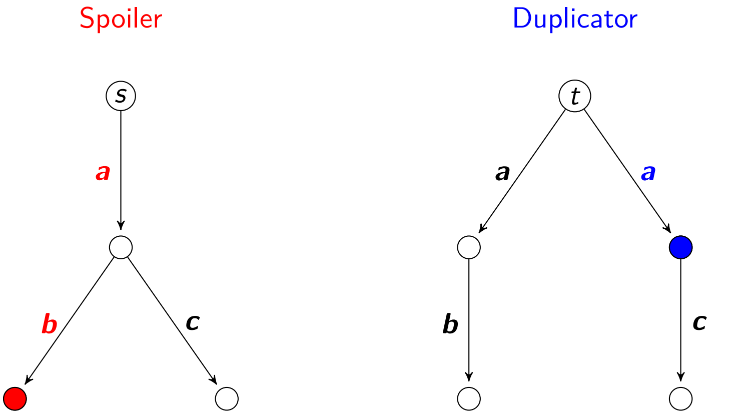
The Simulation Game



The Simulation Game

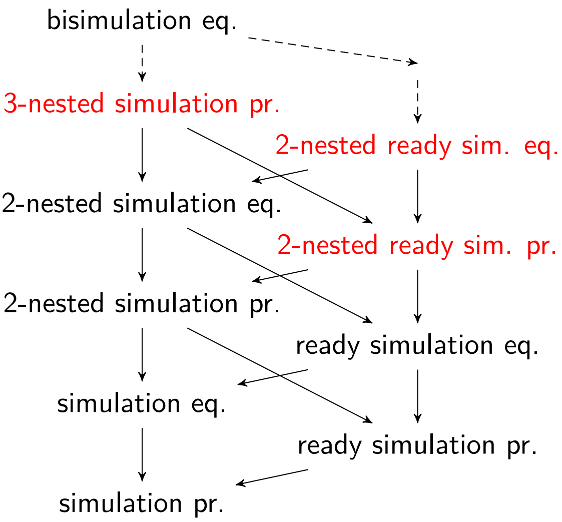


The Simulation Game

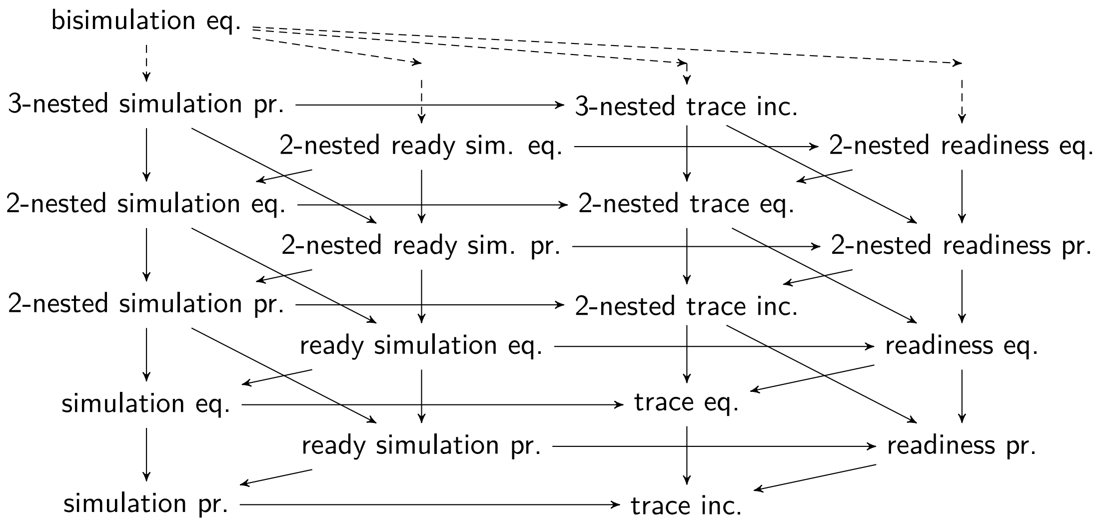


Spoiler wins

The LTBT Spectrum, Game Version



The LTBT Spectrum, Game Version



The Simulation Game, Revisited

1. Player 1 chooses edge from s (leading to s')
2. Player 2 chooses matching edge from t (leading to t')
3. Game continues from configuration s', t'

ω. If Player 2 can always answer: YES, t simulates s.
Otherwise: NO

Or, as an Ehrenfeucht-Fraïssé game ("delayed evaluation"):

1. Player 1 chooses edge from s (leading to s')
2. Player 2 chooses edge from t (leading to t')
3. Game continues from new configuration s', t'

ω. At the end (maybe after infinitely many rounds!), compare the chosen traces:
If the trace chosen by t matches the one chosen by s: YES
Otherwise: NO

Quantitative Ehrenfeucht-Fraïssé Games

The quantitative setting:

- Assume we have a way, possibly application-determined, to measure distances of (finite or infinite) traces
- a hemimetric $D : (\sigma, \tau) \mapsto D(\sigma, \tau) \in \mathbb{R}_{\geq 0} \cup \{\infty\}$

The quantitative Ehrenfeucht-Fraïssé game:

1. Player 1 chooses edge from s (leading to s')
2. Player 2 chooses edge from t (leading to t')
3. Game continues from new configuration s', t'

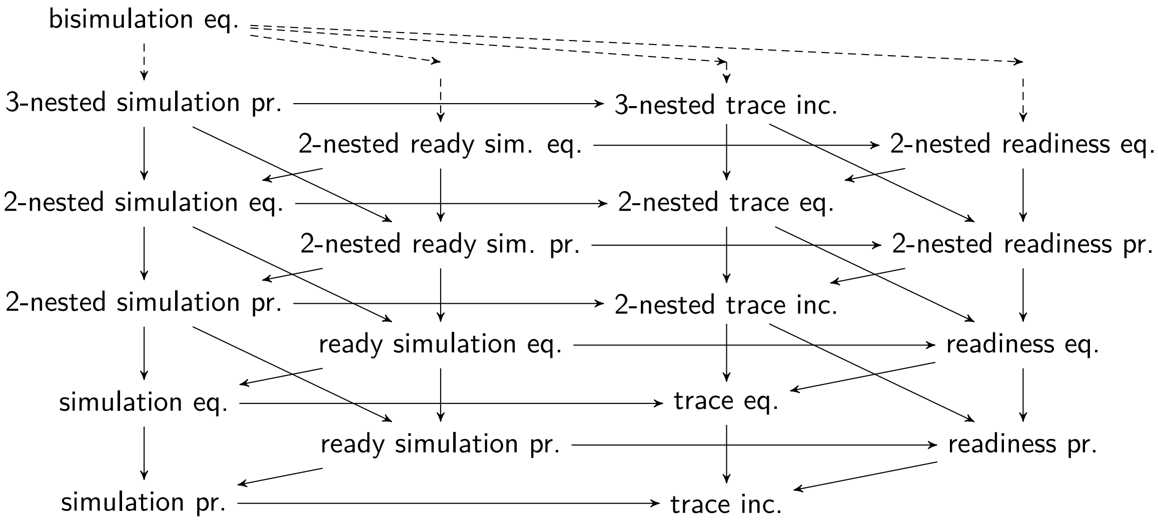
ω. At the end, compare the chosen traces σ, τ:
The simulation distance from s to t is defined to be $D(\sigma, \tau)$

- Player 1 plays to maximize $D(\sigma, \tau)$; Player 2 plays to minimize

This can be generalized to all the games in the LTBT spectrum.

The Quantitative Linear-Time–Branching-Time Spectrum

For any trace distance $D : (\sigma, \tau) \mapsto D(\sigma, \tau) \in \mathbb{R}_{\geq 0} \cup \{\infty\}$:



Quantitative EF Games: Some Details

- Configuration of the game: (π, ρ): π the Player-1 choices up to now; ρ the Player-2 choices
- Strategy: mapping from configurations to next moves
  - Θ_i: set of Player-i strategies
- Simulation strategy: Player-1 moves allowed from end of π
- Bisimulation strategy: Player-1 moves allowed from end of π or end of ρ
  - (hence π and ρ are generally not paths – "mingled paths")
- Pair of strategies $\Longrightarrow$ (possibly infinite) sequence of configurations
- Take the limit; unmingle $\Longrightarrow$ pair of (possibly infinite) traces (σ, τ)
- Bisimulation distance: $\sup_{\theta_1 \in \Theta_1} \inf_{\theta_2 \in \Theta_2} d_T(\sigma, \tau)$
- Simulation distance: $\sup_{\theta_1 \in \Theta_1^0} \inf_{\theta_2 \in \Theta_2} d_T(\sigma, \tau)$ (restricting Player 1's capabilities)

Quantitative EF Games: Some Details – II

- Blind Player-1 strategies: depend only on the end of ρ
  - ("cannot see Player-2 moves")
  - $\tilde\Theta_1$: set of blind Player-1 strategies
- Trace inclusion distance: $\sup_{\theta_1\in\tilde\Theta_1^0} \inf_{\theta_2\in\Theta_2} d_T(\sigma,\tau)$
- For nesting: count the number of times Player 1 switches between end of π and end of ρ
  - Θ_1^k: k switches allowed
- Nested simulation distance: $\sup_{\theta_1\in\Theta_1^1} \inf_{\theta_2\in\Theta_2} d_T(\sigma,\tau)$
- Nested trace inclusion distance: $\sup_{\theta_1\in\tilde\Theta_1^1} \inf_{\theta_2\in\Theta_2} d_T(\sigma,\tau)$ (!)
- For ready: allow extra "I'll see you" Player-1 transition from end of ρ

Transfer Theorem

Theorem

If two equivalences or preorders are inequivalent in the qualitative setting, and the trace distance D is separating, then the corresponding QLTBT distances are topologically inequivalent.

Recursive Characterization

Theorem

If the trace distance $D : (\sigma, \tau) \mapsto d(\sigma, \tau)$ *has a decomposition* $d = g \circ f : \mathrm{Tr} \times \mathrm{Tr} \to L \to \mathbb{R}_{\geq 0} \cup \{\infty\}$ *through a complete lattice* L, *and* f *has a recursive characterization, i.e., such that* $f(a.\sigma, b.\tau) = F(a, b, f(\sigma, \tau))$ *for some* $F : \Sigma \times \Sigma \times L \to L$ *which is monotone in the third coordinate, then all distances in the corresponding QLTBT spectrum are given as least fixed points of some functionals using* F.

All trace distances I know can be expressed recursively like this.

- except ABD'18?
- L is "memory"
- also gives relation family characterization

Specification Theories

Let Mod be a set of models with an equivalence $\sim$.

Definition

A complete specification theory for $(\text{Mod}, \sim)$ is $(\text{Spec}, \leq, \|, \chi)$ such that

- $\leq$ is a refinement preorder on Spec
- $\chi : \text{Mod} \to \text{Spec}$ picks out characteristic specifications
  - *i.e.*, $\forall \mathcal{M}_1, \mathcal{M}_2 \in \text{Mod} : \mathcal{M}_1 \sim \mathcal{M}_2 \iff \chi(\mathcal{M}_1) \leq \chi(\mathcal{M}_2)$
- $(\text{Spec}, \leq, \|)$ forms a bounded commutative distributive residuated lattice up to $\leq \cap \geq$

$\Rightarrow$ $\vee$ and $\wedge$ on Spec; double distributivity; $\bot, \top \in \text{Spec}$

- everything up to modal equivalence $\equiv \; = \; \leq \cap \geq$

$\Rightarrow$ $\|$ distributes over $\vee$, has unit U, has residual $/$ (up to $\equiv$)

- $\mathcal{S}_1 \| \mathcal{S}_2 \leq \mathcal{S}_3 \iff \mathcal{S}_2 \leq \mathcal{S}_3 / \mathcal{S}_1$

Examples

- Disjunctive modal transition systems
- Acceptance automata
- Hennessy-Milner logic with maximal fixed points

- CONCUR 2013, ICTAC 2014, I&C 2020 (all with $\sim =$ bisimulation)

Acceptance Automata

Let Σ be a finite alphabet.

Definition

A (nondeterministic) acceptance automaton (AA) is a structure $\mathcal{A} = (S, S^0, \mathrm{Tran})$, with $S \supseteq S^0$ finite sets of states and initial states and $\mathrm{Tran} : S \to 2^{2^{\Sigma \times S}}$ an assignment of *transition constraints*.

- standard labeled transition system (LTS): $\mathrm{Tran} : S \to 2^{\Sigma \times S}$ (coalgebraic view)
- (for AA:) $\mathrm{Tran}(s) = \{M_1, M_2, \dots\}$: provide M_1 or M_2 or ...
- a disjunctive choice of conjunctive constraints
- J.-B. Raclet 2008 (but deterministic); see also H. H. Hansen 2003
- note multiple initial states

Refinement

Definition

Let $\mathcal{A}_1 = (S_1, S_1^0, \mathrm{Tran}_1)$ and $\mathcal{A}_2 = (S_2, S_2^0, \mathrm{Tran}_2)$ be AA.
A relation $R \subseteq S_1 \times S_2$ is a modal refinement if:

1. $\forall s_1^0 \in S_1^0 : \exists s_2^0 \in S_2^0 : (s_1^0, s_2^0) \in R$ (init)
2. $\forall (s_1, s_2) \in R : \forall M_1 \in \mathrm{Tran}_1(s_1) : \exists M_2 \in \mathrm{Tran}_2(s_2) :$ (tran)
   1. $\forall (a, t_1) \in M_1 : \exists (a, t_2) \in M_2 : (t_1, t_2) \in R$
   2. $\forall (a, t_2) \in M_2 : \exists (a, t_1) \in M_1 : (t_1, t_2) \in R$

Write $\mathcal{A}_1 \le \mathcal{A}_2$ if there exists such a modal refinement.

- for any constraint choice M_1 there is a bisimilar choice M_2
- $\mathcal{A}_1$ has fewer choices than $\mathcal{A}_2$
- no more choices $\triangleq$ only one $M \in \mathrm{Tran}(s)$ $\triangleq$ LTS
- formally: an embedding $\chi : \mathrm{LTS} \hookrightarrow \mathrm{AA}$
  such that $\chi(\mathcal{L}_1) \le \chi(\mathcal{L}_2)$ iff $\mathcal{L}_1$ and $\mathcal{L}_2$ are bisimilar

Logical Operations

Let $\mathcal{A}_1 = (S_1, S_1^0, \text{Tran}_1)$ and $\mathcal{A}_2 = (S_2, S_2^0, \text{Tran}_2)$ be AA.

Disjunction: $\mathcal{A}_1 \vee \mathcal{A}_2 = (S_1 \overset{+}{\cup} S_2, S_1^0 \overset{+}{\cup} S_2^0, \text{Tran}_1 \overset{+}{\cup} \text{Tran}_2)$

Conjunction: define $\pi_i : 2^{\Sigma \times S_1 \times S_2} \to 2^{\Sigma \times S_i}$ by

$$\pi_1(M) = \{(a, s_1) \mid \exists s_2 \in S_2 : (a, s_1, s_2) \in M\}$$
$$\pi_2(M) = \{(a, s_2) \mid \exists s_1 \in S_1 : (a, s_1, s_2) \in M\}$$

Let $\mathcal{A}_1 \wedge \mathcal{A}_2 = (S_1 \times S_2, S_1^0 \times S_2^0, \text{Tran})$ with

$$\text{Tran}((s_1, s_2)) = \{M \subseteq \Sigma \times S_1 \times S_2 \mid \pi_1(M) \in \text{Tran}_1(s_1), \pi_2(M) \in \text{Tran}_2(s_2)\}$$

Theorem

For all LTS $\mathcal{L}$ and AA $\mathcal{A}_1, \mathcal{A}_2$:

$$\mathcal{L} \models \mathcal{A}_1 \vee \mathcal{A}_2 \iff \mathcal{L} \models \mathcal{A}_1 \text{ or } \mathcal{L} \models \mathcal{A}_2$$
$$\mathcal{L} \models \mathcal{A}_1 \wedge \mathcal{A}_2 \iff \mathcal{L} \models \mathcal{A}_1 \ \& \ \mathcal{L} \models \mathcal{A}_2$$

Structural Operations: Composition

Let $\mathcal{A}_1 = (S_1, S_1^0, \mathrm{Tran}_1)$ and $\mathcal{A}_2 = (S_2, S_2^0, \mathrm{Tran}_2)$ be AA.

For $M_1 \subseteq \Sigma \times S_1$ and $M_2 \subseteq \Sigma \times S_2$, define

$$M_1 \| M_2 = \{(a, (t_1, t_2)) \mid (a, t_1) \in M_1, (a, t_2) \in M_2\}$$

(assumes CSP synchronization, but can be generalized)

Let $\mathcal{A}_1 \| \mathcal{A}_2 = (S_1 \times S_2, S_1^0 \times S_2^0, \mathrm{Tran})$ with

$$\mathrm{Tran}((s_1, s_2)) = \{M_1 \| M_2 \mid M_1 \in \mathrm{Tran}_1(s_1), M_2 \in \mathrm{Tran}_2(s_2)\}$$

Theorem (independent implementability)

For all AA $\mathcal{A}_1$, $\mathcal{A}_2$, $\mathcal{A}_3$, $\mathcal{A}_4$:

$$\mathcal{A}_1 \le \mathcal{A}_3 \;\&\; \mathcal{A}_2 \le \mathcal{A}_4 \implies \mathcal{A}_1 \| \mathcal{A}_2 \le \mathcal{A}_3 \| \mathcal{A}_4$$

Structural Operations: Quotient

Let $\mathcal{A}_1 = (S_1, S_1^0, \text{Tran}_1)$ and $\mathcal{A}_2 = (S_2, S_2^0, \text{Tran}_2)$ be AA.

Define $\mathcal{A}_1 / \mathcal{A}_2 = (S, S^0, \text{Tran})$:

- $S = 2^{S_1 \times S_2}$
- write $S_2^0 = \{s_2^{0,1}, \dots, s_2^{0,p}\}$ and let $S^0 = \{\{(s_1^{0,q}, s_2^{0,q}) \mid q \in \{1, \dots, p\}\} \mid \forall q : s_1^{0,q} \in S_1^0\}$
- $\text{Tran} =$

Structural Operations: Quotient

Let $\mathcal{A}_1 = (S_1, S_1^0, \mathrm{Tran}_1)$ and $\mathcal{A}_2 = (S_2, S_2^0, \mathrm{Tran}_2)$ be AA.

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- $\mathrm{Tran} =$

Structural Operations: Quotient

Let $\mathcal{A}_1 = (S_1, S_1^0, \mathrm{Tran}_1)$ and $\mathcal{A}_2 = (S_2, S_2^0, \mathrm{Tran}_2)$ be AA.

Define $\mathcal{A}_1 / \mathcal{A}_2 = (S, S^0, \mathrm{Tran})$:

- $S = 2^{S_1 \times S_2}$
- write $S_2^0 = \{s_2^{0,1}, \dots, s_2^{0,p}\}$ and let $S^0 = \{\{(s_1^{0,q}, s_2^{0,q}) \mid q \in \{1, \dots, p\}\} \mid \forall q : s_1^{0,q} \in S_1^0\}$
- $\mathrm{Tran} = \quad \dots$

Theorem

For all AA $\mathcal{A}_1$, $\mathcal{A}_2$, $\mathcal{A}_3$:

$$\mathcal{A}_1 \| \mathcal{A}_2 \le \mathcal{A}_3 \iff \mathcal{A}_2 \le \mathcal{A}_3 / \mathcal{A}_1$$

- up to $\equiv$, $/$ is the adjoint (or residual) of $\|$

Quantitative Specification Theories?

Definition (recall)

A complete specification theory for $(\mathsf{Mod}, \sim)$ is $(\mathsf{Spec}, \leq, \|, \chi)$ such that

- $\leq$ is a refinement preorder on Spec
- $\mathcal{M}_1 \sim \mathcal{M}_2 \iff \chi(\mathcal{M}_1) \leq \chi(\mathcal{M}_2)$
- $(\mathsf{Spec}, \leq, \|)$ forms a b.c.d. residuated lattice up to $\equiv$

- generalize $\sim$ by pseudometric d_{Mod}
  - $d_{\mathsf{Mod}}(\mathcal{M}_1, \mathcal{M}_2) = 0$ iff $\mathcal{M}_1 \sim \mathcal{M}_2$
- generalize $\leq$ by hemimetric d
  - $d_{\mathsf{Mod}}(\mathcal{M}_1, \mathcal{M}_2) = d(\chi(\mathcal{M}_1), \chi(\mathcal{M}_2))$
  - $d(\mathcal{M}, \mathcal{S}) = d(\chi(\mathcal{M}), \mathcal{S})$
- still want $(\mathsf{Spec}, \leq, \|)$ to be a b.c.d. residuated lattice up to $\equiv$

Acceptance Automata

For DMTS/AA/HML_{max}:

- d_{Mod}: any bisimulation distance
- d: corresponding modal refinement distance
- transitivity $\rightsquigarrow$ triangle ineq.: $d(\mathcal{S}_1, \mathcal{S}_2) + d(\mathcal{S}_2, \mathcal{S}_3) \geq d(\mathcal{S}_1, \mathcal{S}_3)$
- $d(\mathcal{S}, \mathcal{S}_1 \wedge \mathcal{S}_2) = \max(d(\mathcal{S}, \mathcal{S}_1), d(\mathcal{S}, \mathcal{S}_2))$ or ∞
- $d(\mathcal{S}_1 \vee \mathcal{S}_2, \mathcal{S}) = \max(d(\mathcal{S}_1, \mathcal{S}), d(\mathcal{S}_2, \mathcal{S}))$ or ∞
- quotient is quantitative residual: $d(\mathcal{S}_1 \| \mathcal{S}_2, \mathcal{S}_3) = d(\mathcal{S}_2, \mathcal{S}_3 / \mathcal{S}_1)$
- for $\|$ itself, uniform continuity: a function $P : \mathbb{R}_{\geq 0} \times \mathbb{R}_{\geq 0} \to \mathbb{R}_{\geq 0}$ such that $d(\mathcal{S}_1 \| \mathcal{S}_2, \mathcal{S}_3 \| \mathcal{S}_4) \leq P(d(\mathcal{S}_1, \mathcal{S}_3), d(\mathcal{S}_2, \mathcal{S}_4))$

The Bad
and/or
Ugly

Silent Moves in QLTBT?

- Any serious spectrographer needs to think about silent moves
- (van Glabbeek 1993: LTBT II)
- Bisping, Jansen 2023: Energy games for the weak spectrum
  - but uses power set for linear part (recall: we use blindness instead)
  - difficult to reconcile power set with quantitative setting
- otherwise, some coalgebra approaches:
  - Sprunger, Katsumata, Dubut, Hasuo 2021: Fibrational bisimulations and quantitative reasoning
  - Ford, Milius, Schröder, Beohar, König 2022: Graded monads and behavioural equivalence games
  - Beohar, Gurke, König, Messing 2023: Hennessy-Milner theorems via Galois connections
  - again, power set seems very popular ...
- status: IT'S COMPLICATED

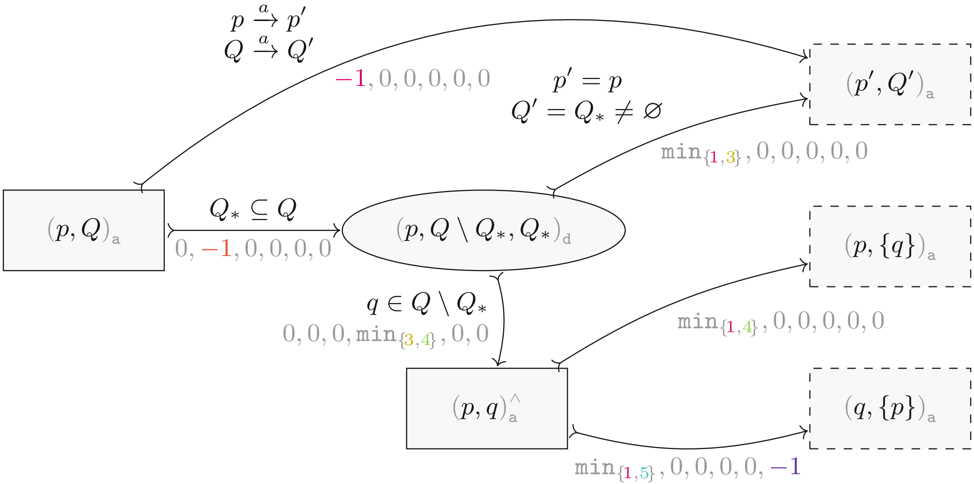


Fig. 7. Schematic spectroscopy game $\mathcal{G}_\triangle$ of Definition 10.

Asarin-Basset-Degorre Distance

Recall:

$$D(\sigma, \tau) = \max \begin{cases} \sup_i \inf_j \{|t_i - s_j| \mid a_i = b_j\} \\ \sup_j \inf_i \{|t_i - s_j| \mid a_i = b_j\} \end{cases}$$

- takes into account permutations of symbols which are close in timing
- but in a way which may lose symbols
- relation to timed pomsets? Amrane, Bazille, Clement, UF 2024: Languages of HDTA
- status: HOPEFUL

On the practical side, if we observed timed words with some finite precision (say $0.01s$), then it would be difficult to distinguish the order of close events, e.g. detect the difference between

$$w_1 = (a, 1), (b, 2), (c, 2.001) \text{ and } w_2 = (a, 1.001), (c, 1.999), (b, 2.001).$$

Moreover, it is even difficult to count the number of events that happen in a short lapse of time, e.g. the words w_1, w_2 look very similar to

$$w_3 = (a, 1), (c, 1.999), (c, 2), (b, 2.001), (c, 2.0002).$$

A slow observer, when receiving timed words w_1, w_2, w_3 will just sense an a at the date ≈ 1 and b and c at the date ≈ 2.

As the main contribution of this paper, we introduce a metric on timed words (with non-fixed number of events) for which w_1, w_2, w_3 are very close to each other. We believe that this metric is natural and sets a ground for approximate model-checking and information theory of timed languages w.r.t. time (and not only number of events).

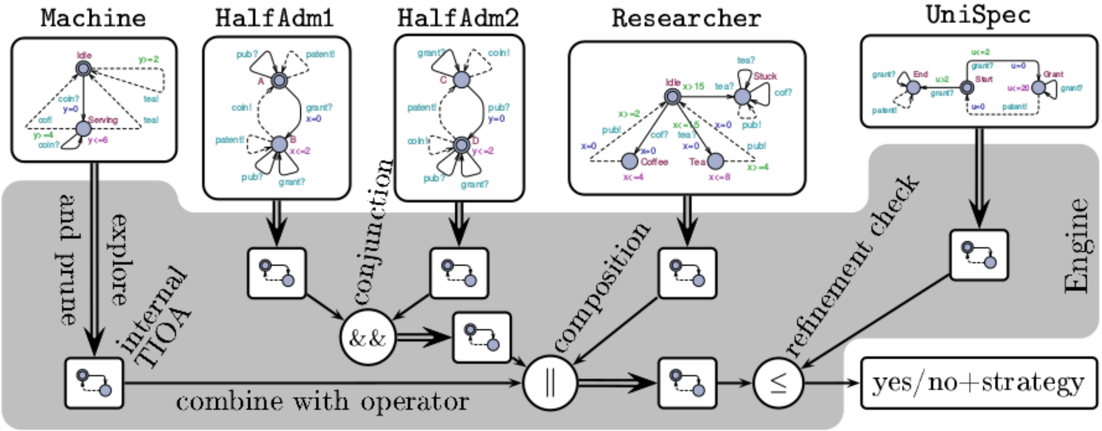
We present the first technical results concerning this distance:

Specification Theories for Real-Time Systems

Timed input-output automata:

- David, Larsen, Legay, Nyman, Traonouez, Wąsowski 2015: Real-time specifications
- Goorden, Larsen, Legay, Lorber, Nyman, Wąsowski 2023: Timed I/O Automata: It is never too late to complete your timed specification theory
- complete, with quotient, but without disjunction
- only deterministic specifications
- tool support: ECDAR / Uppaal TiGa (Aalborg)
- some work on robustness and implementability: Larsen, Legay, Traonouez, Wąsowski 2014: Robust synthesis for real-time systems

Timed Input-Output Automata



Specification Theories for Real-Time Systems, contd.

Modal event-clock specifications:

- Bertrand, Legay, Pinchinat, Raclet 2012: Modal event-clock specifications for timed component-based design
- complete, with quotient, but without disjunction
- only deterministic specifications
- some work on robustness: UF, Legay 2012: A robust specification theory for modal event-clock automata

Synchronous time-triggered interface theories:

- Delahaye, UF, Henzinger, Legay, Ničković 2012: Synchronous interface theories and time triggered scheduling
- no quotient, dubious conjunction, no implementation
- relation to BIP

Specification Theories for Hybrid Systems

Specification Theories for Hybrid Systems

- Quesel, Fränzle, Damm 2011: Crossing the bridge between similar games

Conclusion

- general theory of quantitative verification ✓
- general theory of compositional quantitative verification ¯_(ツ)_/¯
  - algebraic properties ✓
  - quantitative algebraic properties ✗
  - silent moves ✗
- for real-time systems ¯_(ツ)_/¯
  - robustness ¯_(ツ)_/¯
  - compositionality ¯_(ツ)_/¯
  - robust compositionality ✗
- for hybrid systems ✗